

# Digit expansions in Rational and Algebraic Basis

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# RATIONAL BASE NUMBER SYSTEMS

(Akiyama, Frougny and Sakarovitch, '08)

$$\frac{a}{b} \in \mathbb{Q}, \quad a > b \geq 2 \quad \text{base}$$

$$\mathcal{D} = \{0, 1, \dots, a-1\} \quad \text{digit set}$$

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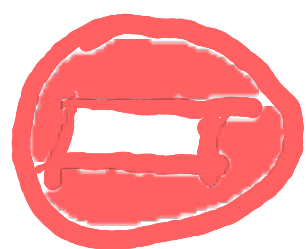
Given  $N \in \mathbb{N}$ , they introduced an expansion

$$N = \frac{1}{b} \sum_{j=0}^k \left(\frac{a}{b}\right)^j d_j$$

( $k \geq 0$ ,  $d_j \in \mathcal{D}$ ) through the modified division algorithm.

$$N = N_0, \quad j \geq 0, \quad \boxed{b N_j = a N_{j+1} + d_j}$$

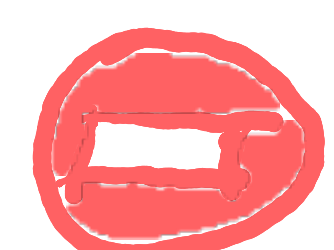
where  $d_j \equiv b N_j \pmod{a}$  (recall  $\mathcal{D} = \{0, \dots, a-1\}$ ).

Then  $N_{k+1} = 0$  for some  $k$ . 



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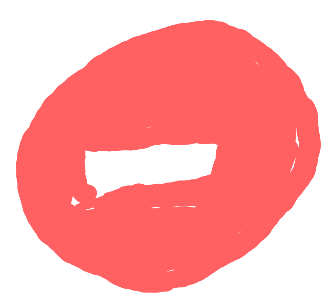
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Example:  $\frac{a}{b} = \frac{3}{2}$ ,  $\mathcal{D} = \{0, 1, 2\}$ ,  $N = 4$

$$2 \cdot \underline{4} = 3 \cdot N_1 + d_0 \implies d_0 = 2 \quad N_1 = 2$$

$$2 \cdot 2 = 3 \cdot N_2 + d_1 \implies d_1 = 1 \quad N_2 = 1$$

$$2 \cdot 1 = 3 \cdot N_3 + d_2 \implies d_2 = 2 \quad \underline{N_3 = 0} \quad \text{$$

$$\begin{aligned} \underline{4} &= \frac{3}{2} N_1 + \frac{2}{2} = \left(\frac{3}{2}\right)^2 N_2 + \frac{3}{2} \cdot \frac{1}{2} + \frac{2}{2} \\ &= \frac{1}{2} \left( \left(\frac{3}{2}\right)^2 \cdot 2 + \left(\frac{3}{2}\right) \cdot 1 + 2 \right) = (212)_{\frac{3}{2}} \end{aligned}$$

Reinterpreting the algorithm:

$$N_j \in \mathbb{N}$$

$$b N_j = a N_{j+1} + d_j$$

Now:  $\frac{a}{b}$  can be negative,

$$N_j \in b\mathbb{Z}$$

$$N_j = \frac{a}{b} \cdot N_{j+1} + d_j$$

$$\Rightarrow N = \left(\frac{a}{b}\right)^k d_k + \dots + \left(\frac{a}{b}\right) d_1 + d_0$$

(I can get rid of the  $\frac{1}{b}$  factor)

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$$T_{a/b}: \begin{array}{ccc} b\mathbb{Z} & \longrightarrow & b\mathbb{Z} \\ N & \longmapsto & \frac{N-d}{a/b} \end{array} \quad d \in \mathbb{D}$$

backward division map restricted to a lattice

# ALGEBRAIC NUMBER SYSTEMS

$$\alpha \in \mathbb{Q}(i), \quad |\alpha| > 1 \quad \alpha \notin \mathbb{R}$$

$$P_\alpha(X) = a_2 X^2 + a_1 X + a_0 \quad \text{minimal polynomial}$$

$$a_2 \neq 1$$

$$\mathcal{D} = \{0, 1, \dots, |a_0| - 1\}$$

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We need a lattice  $\Lambda_\alpha \subset \mathbb{C}$  such that

$$\begin{array}{ccc} T_\alpha: \Lambda_\alpha & \longrightarrow & \Lambda_\alpha \\ x & \longmapsto & \frac{x - d}{\alpha} \end{array}$$

is well defined,

where  $d$  is the unique digit in  $\mathcal{D}$  such that

$$\frac{x - d}{\alpha} \in \Lambda_\alpha.$$



$$\Lambda_\alpha := \mathbb{Z}[\alpha] \cap \alpha^{-1} \mathbb{Z}[\alpha^{-1}]$$

Brunotte basis:  $\Lambda_\alpha = a_2 \mathbb{Z} + (a_2 \alpha + a_1) \mathbb{Z}$

$$(a_2 \alpha^2 + a_1 \alpha + a_0 = 0)$$



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 $(a_2 \alpha^2 + a_1 \alpha + a_0 = 0)$

Algorithm:

$$N = N_0 \in \Lambda_\alpha$$

$$N_j = \alpha N_{j+1} + d_j$$

where  $d_j$  is the unique digit

such that  $N_{j+1} \in \Lambda_\alpha$

Example:  $\alpha = \frac{-1+3i}{2}$        $P_\alpha(x) = 2x^2 + 2x + 5$

$$\mathcal{D} = \{0, 1, 2, 3, 4\}$$

$$\Lambda_\alpha = a_2 \mathbb{Z} + (a_2 \alpha + a_1) \mathbb{Z} = 2\mathbb{Z} + (1+3i)\mathbb{Z}$$

Let  $N = \underline{1+3i} \in \Lambda_\alpha$

$$\underline{1+3i} = \frac{-1+3i}{2} \cdot 2 + 2$$

$$2 = \frac{-1+3i}{2} \cdot 0 + 2$$

$$\Rightarrow \underline{1+3i} = (2, 2)_\alpha$$

Does the algorithm always terminate?  
that is, is there always  $k \geq 0$  s.t.  $N_k = 0$ ?

No, and the set of bases  $\alpha$  with this property is extremely difficult to describe.

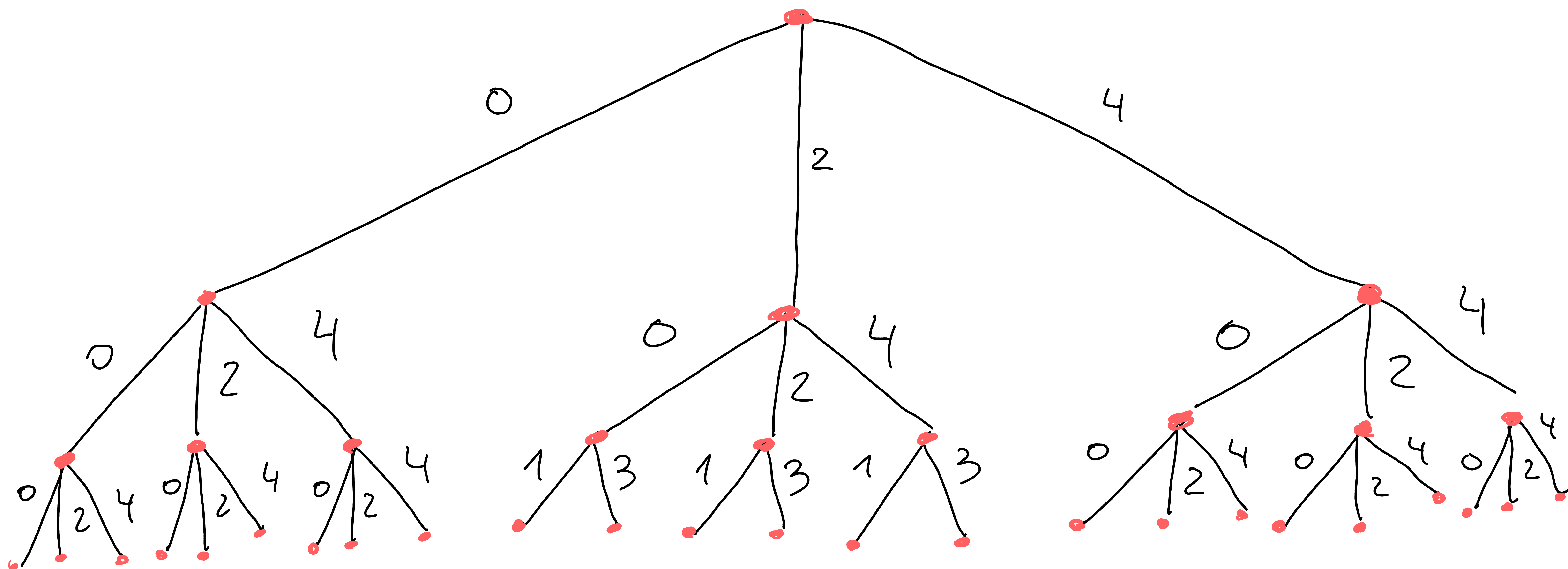
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This set has been well studied in the context of **shift radix systems**.

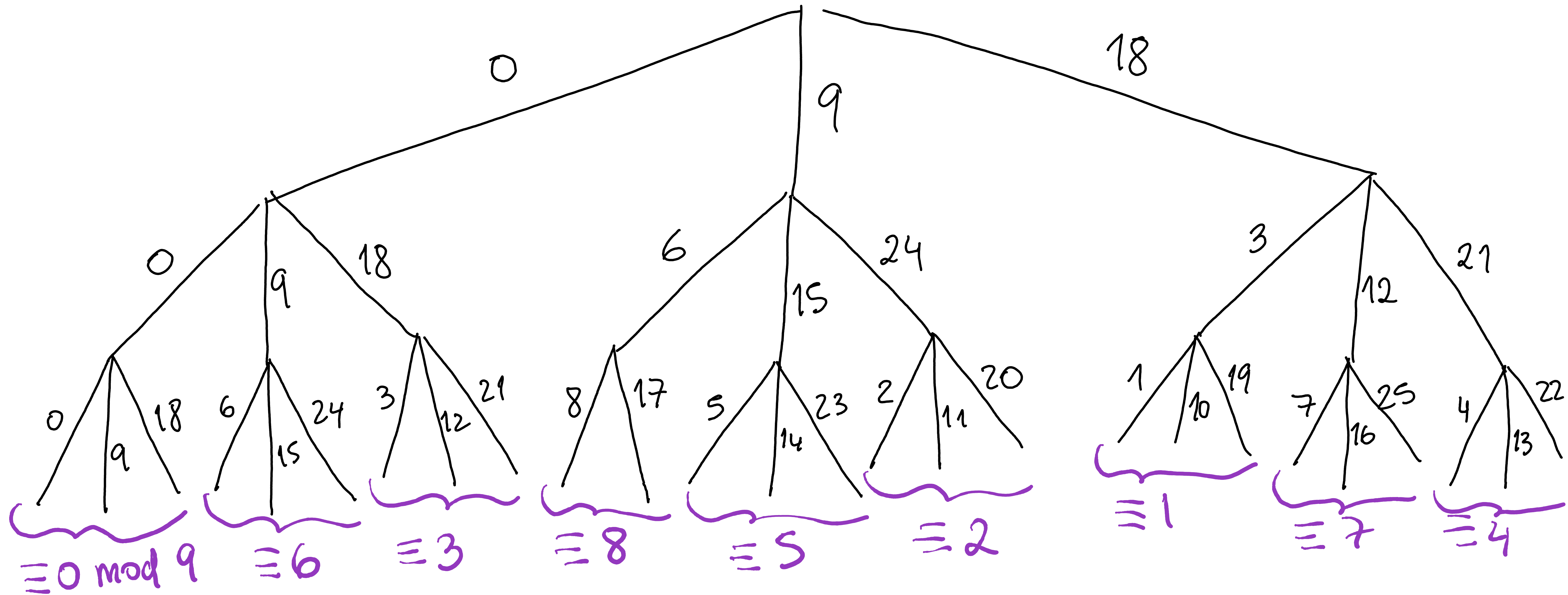
We assume from now on that  $\alpha$  has the **finiteness property**.

The tree of  $\alpha$ -expansions for  $\alpha = \frac{-1+3i}{2}$





The tree for  $\alpha = \frac{-1+5i}{3}$ ,  $P_\alpha = 9X^2 + 6X + 26$   
 $\mathcal{D} = \{0, 1, \dots, 25\}$





$L_\alpha$  = language of expansions in base  $\alpha$

$L_\alpha^k$  = expansions in base  $\alpha$  of length  $k$

Proposition:

$$w \in L_\alpha^k$$

$$\{d \in \mathcal{D} : wd \in L_\alpha^{k+1}\} = \{d \in \mathcal{D} : d \equiv r \pmod{a_2}\}$$

for some  $r$ .

Corollary:

$$|L_\alpha^k| \leq \left\lceil \frac{|\mathcal{D}|}{a_2} \right\rceil^k$$

$$\forall k \geq 0.$$

# EXPANSION OF COMPLEX NUMBERS

Def: An  $\alpha$ -expansion of  $x \in \mathbb{C}$  is an expansion of the form

$$x = (d_k \dots d_0 \cdot d_{-1} d_{-2} \dots)_\alpha = \sum_{j \leq k} d_j \alpha^j$$

such that each finite prefix

$$d_k \dots d_\ell$$

is the  $\alpha$ -expansion of some  $N \in \Lambda_\alpha$

(i.e., infinite paths on the tree).

# FRACTAL TILES

$$z \in \Lambda_\alpha$$

$$z = (d_k \dots d_0)_\alpha$$

$$g(z) := \{x \in \mathbb{C} : x = \overbrace{(d_k \dots d_0)} \cdot d_{-1} d_{-2} \dots\}_\alpha\}$$

Expansions whose integer part is  $z$ .

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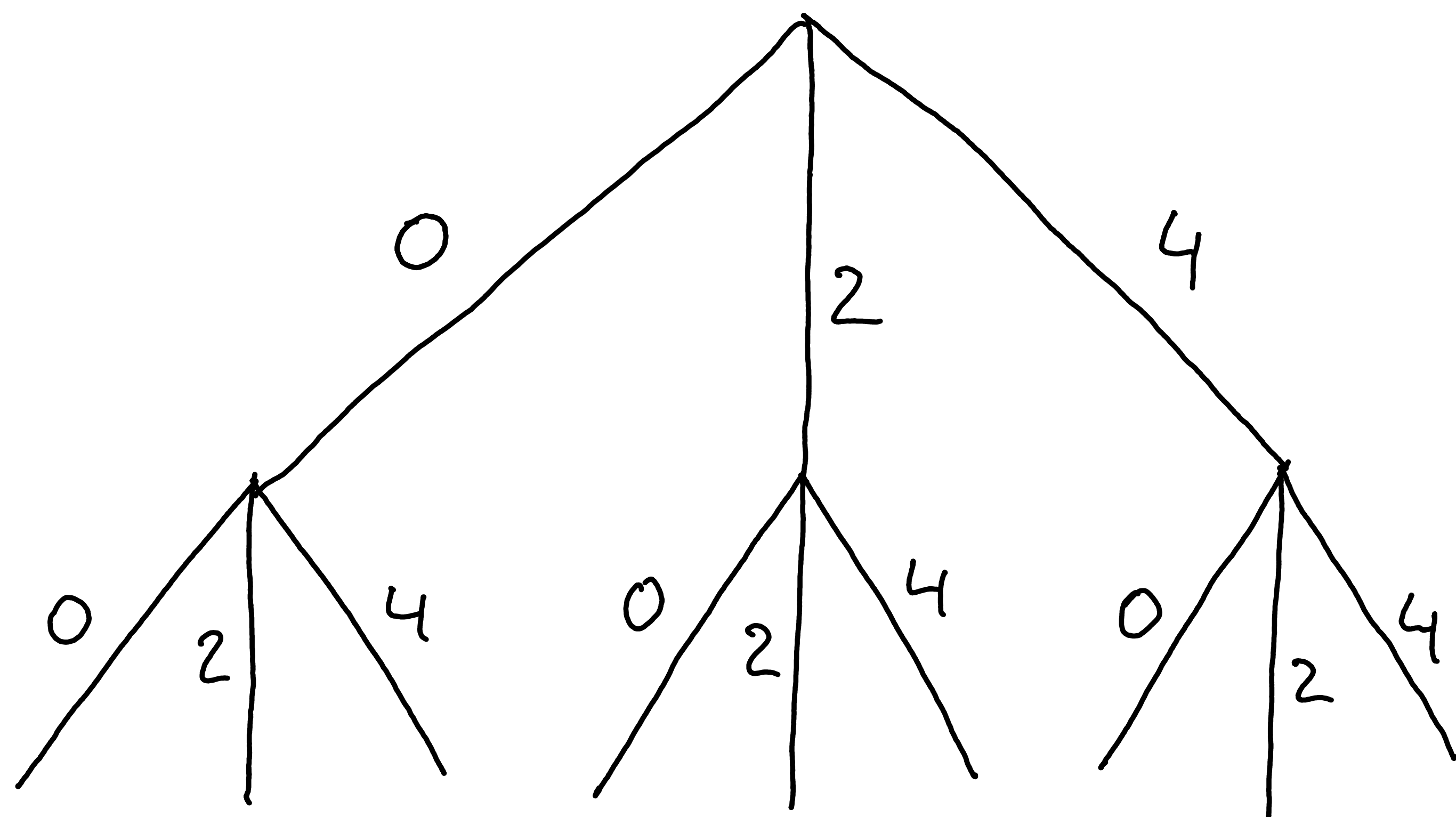
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$$\text{Example: } \alpha = \frac{-1+3i}{2}$$

$$\sqrt{2} = (2.2341121\dots)_\alpha \in g(2)$$

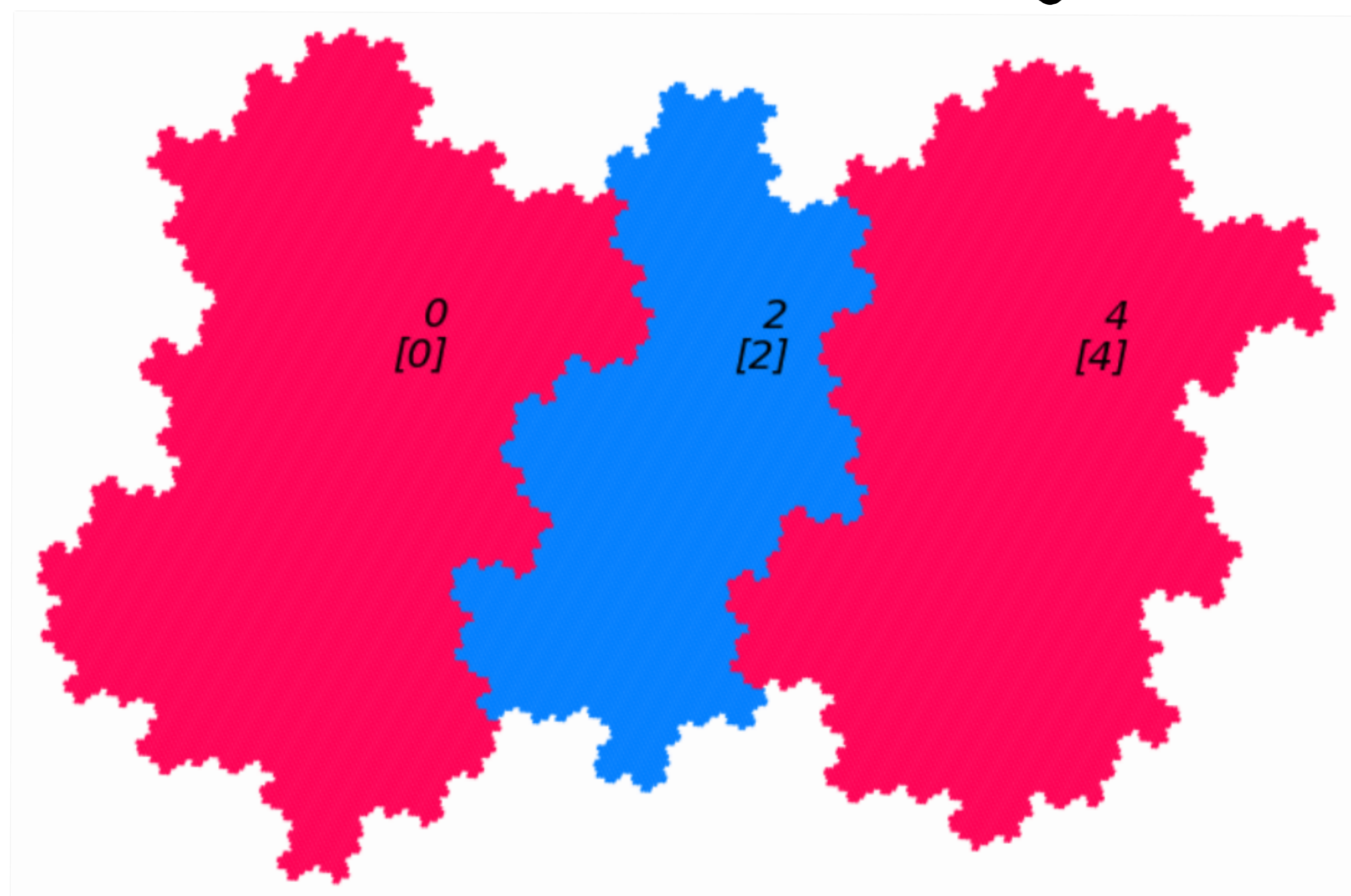




$g(0)$

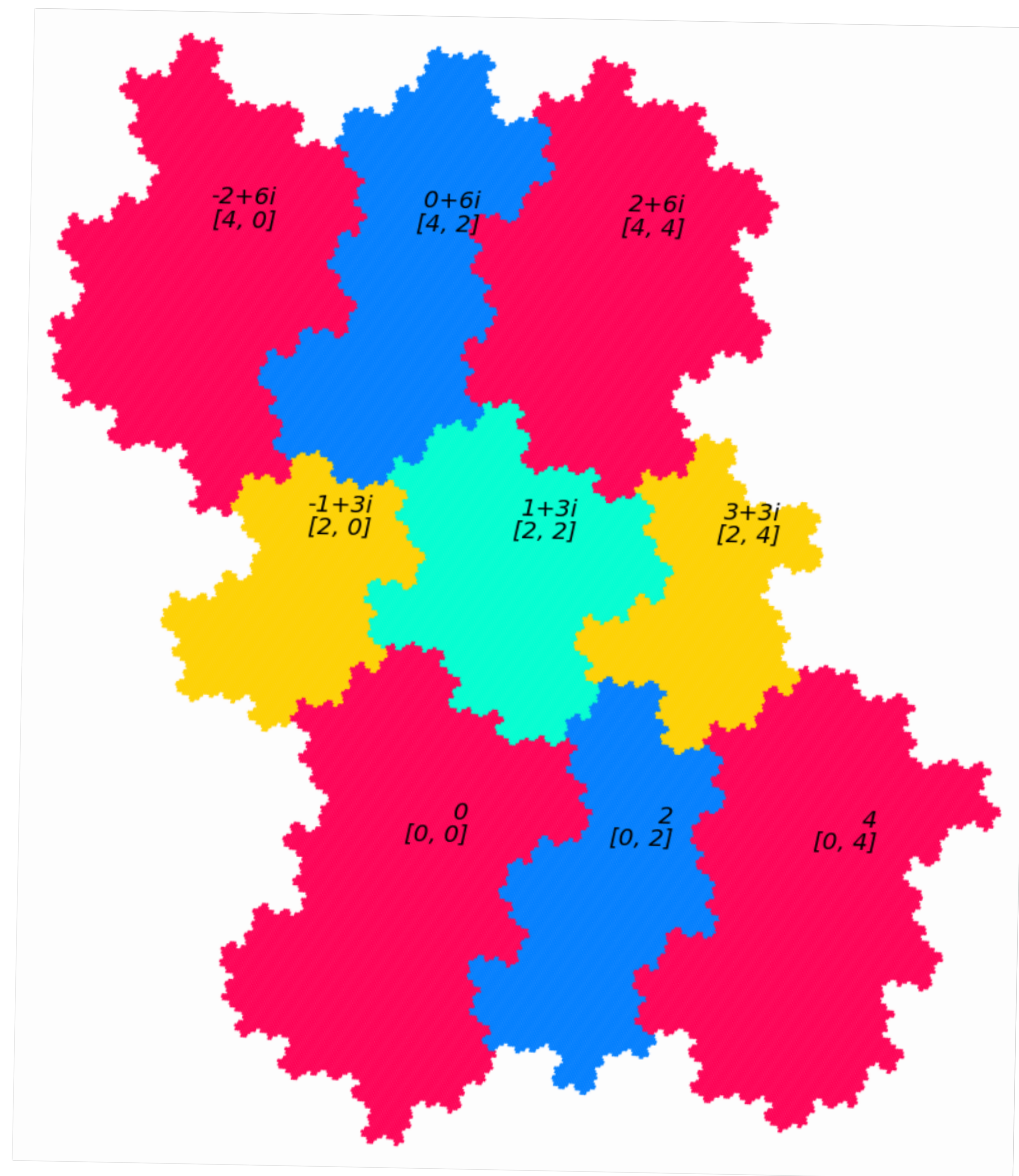
$g(2)$

$g(4)$

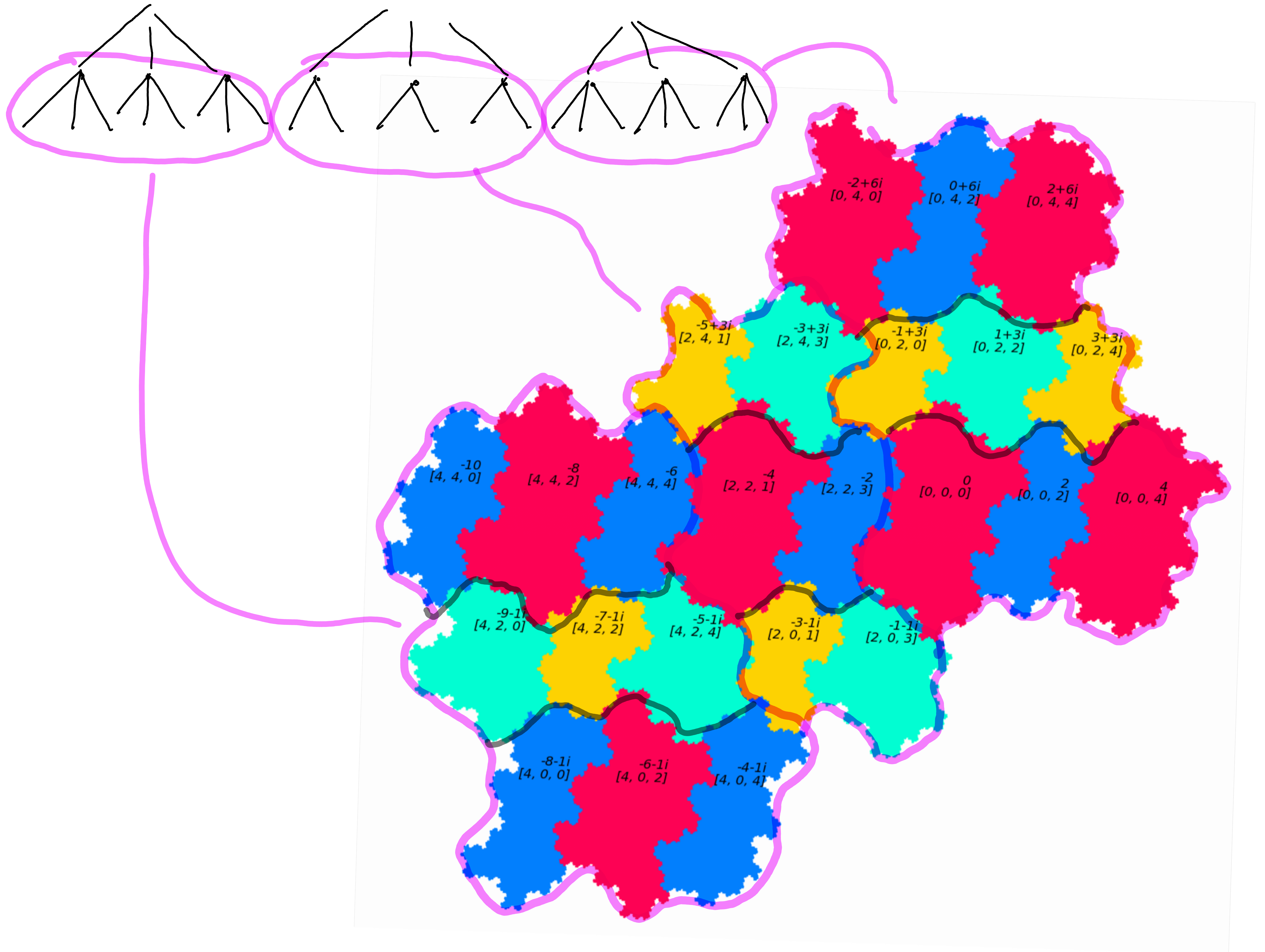


$$\alpha g(0) = g(0) \cup g(2) \cup g(4)$$

$$\alpha = \frac{-1+3i}{2}, \{0,1,2,3,4\}$$









# P-ADIC COMPLETIONS

Steiner, Thurwaldner. "Rational self-affine tiles"

$p \in \mathbb{Q}(i)$  Gaussian prime

$K_p$   $p$ -adic completion of  $\mathbb{Q}(i)$

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$\alpha \in \mathbb{Q}(i) \rightarrow$  unique factorization into Gaussian primes

$$\alpha = \frac{\text{num}(\alpha) \in \mathbb{Z}[i]}{\text{den}(\alpha) \in \mathbb{Z}[i]} \quad (\text{coprime})$$

$$\text{den}(\alpha) = p_1^{r_1} \cdot \dots \cdot p_s^{r_s}$$

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Theorem: a series  $\sum_{j \leq k} d_j \alpha^j$  is an  $\alpha$ -expansion if and only if it converges to 0 in  $K_p \quad \forall p \mid \text{den}(\alpha)$

Theorem:

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Corollary:

The  $\alpha$ -expansion of a complex number is unique almost everywhere.

# OPEN QUESTIONS

- What are exactly the points with multiple  $\alpha$ -expansions?



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- What are exactly the points with multiple  $\alpha$ -expansions?

- Let  $L_\alpha^w$  be the language of all (infinite)  $\alpha$ -expansions.

Do all digits (resp. finite words) appear with the same frequency?

THANK  
you!